IN THE CLAIMS:

- 1. (Currently Amended) A method for a storage operating system implemented in a
- storage system to optimize the amount of readahead data retrieved for a read stream es-
- tablished in a data container stored in the storage system, the method comprising:
- 4 receiving a client read request at the storage system, the client read request indi-
- 5 cating client-requested data for the storage operating system to retrieve from the data
- 6 container containing the read stream;
- determining whether the storage operating system is permitted to retrieve reada-
- head data from the data container in response to the received client read request;
- if it is determined that the storage operating system is permitted to retrieve reada-
- head data from the data container, performing the steps of:
 - (i) selecting an amount of readahead data to retrieve from the data container based on a plurality of factors stored within a readset data structure associated with the read stream; and
 - (ii) retrieving the selected amount of readahead data from the data container.
- 1 2. (Original) The method of claim 1, wherein the data container is a file, directory, vdisk
- 2 or lun.

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- 3. (Original) The method of claim 1, wherein the storage operating system is determined
- to be permitted to retrieve readahead data from the data container when the client-
- requested data extends the read stream past a predetermined next readahead value.
- 4. (Original) The method of claim 3, wherein the predetermined next readahead value is
- stored in a readset data structure associated with the read stream.

- 5. (Original) The method of claim 3, wherein the predetermined next readahead value is
- 2 updated based on a percentage of the selected amount of readahead data.
- 6. (Previously Presented) The method of claim 1, wherein a read-access style associated
- with the data container is one of the plurality of factors used to select the amount of
- 3 readahead data.
- 7. (Original) The method of claim 6, wherein the selected amount of readahead data
- equals zero if the read-access style corresponds to a random read-access style.
- 8. (Previously Presented) The method of claim 1, wherein a number of client read re-
- quests processed in the read stream is one of the plurality of factors used to select the
- 3 amount of readahead data.
- 9. (Original) The method of claim 8, wherein the number of client read requests proc-
- 2 essed in the read stream is stored as a count value in a readset data structure associated
- with the read stream.
- 1 10. (Previously Presented) The method of claim 1, wherein the amount of client-
- 2 requested data is one of the plurality of factors used to select the amount of readahead
- 3 data.
- 1 11. (Original) The method of claim 10, wherein the selected amount of readahead data is
- set equal to a predetermined upper limit for large amounts of client-requested data.
- 1 12. (Original) The method of claim 1, wherein the selected amount of readahead data is
- doubled if the number of client read requests processed in the read stream is greater than
- 3 a first threshold value.

- 1 13. (Original) The method of claim 1, wherein the client-requested data is identified as
- read-once data when either (i) the number of client read requests processed in the read
- stream is greater than a second threshold value or (ii) a set of metadata associated with
- 4 the read stream indicates that the client-requested data is read-once data.
- 14. (Original) The method of claim 1, wherein the selected amount of readahead data is
- stored in one or more buffers enqueued on a flush queue, the flush queue being config-
 - 3 ured to reuse buffers after a predetermined period of time.
 - 15. (Original) The method of claim 14, wherein the predetermined period of time equals
 - 2 two seconds.
 - 1 16. (Currently Amended) An apparatus configured to implement a storage operating sys-
 - tem that optimizes the amount of readahead data retrieved for a read stream established in
 - a data container stored in the apparatus, the apparatus comprising:
 - 4 means for receiving a client read request, the client read request indicating client-
 - 5 requested data for the storage operating system to retrieve from the data container con-
 - 6 taining the read stream;
 - 7 means for determining whether the storage operating system is permitted to re-
 - 8 trieve readahead data from the data container in response to the received client read re-
 - 9 quest;
 - means for selecting an amount of readahead data to retrieve from the data con-
- tainer based on a plurality of factors stored within a readset data structure associated with
- 12 the read stream; and
- means for retrieving the selected amount of readahead data from the data con-
- 14 tainer.
- 17. (Original) The apparatus of claim 16, wherein the data container is a file, directory,
- 2 vdisk or lun.

- 18. (Original) The apparatus of claim 16, wherein the storage operating system is deter-
- 2 mined to be permitted to retrieve readahead data from the data container when the client-
- requested data extends the read stream past a predetermined next readahead value.
- 1 19. (Original) The apparatus of claim 18, further comprising means for updating the pre-
- determined next readahead value based on a percentage of the selected amount of reada-
- 3 head data.
- 20. (Previously Presented) The apparatus of claim 16, wherein the plurality of factors
- 2 used to select the amount of readahead data includes at least one of:
- 3 (i) the amount of client-requested data,
- 4 (ii) a number of client read requests processed in the read stream, and
- 5 (iii) a read-access style associated with the data container.
- 1 21. (Original) The apparatus of claim 16, wherein the selected amount of readahead data
- 2 is doubled if the number of client read requests processed in the read stream is greater
- 3 than a first threshold value.
- 22. (Currently Amended): A storage system configured to optimize the amount of reada-
- 2 head data retrieved for a read stream established in a data container stored in the storage
- 3 system, the storage system comprising:
- a network adapter for receiving a client read request, the client read request indi-
- 5 cating client-requested data to retrieve from the data container containing the read stream;
- 6 and
- a memory configured to store instructions for implementing a storage operating
- system that performs the steps of:
- determining whether the storage operating system is permitted to retrieve
- readahead data from the data container in response to the received client read re-
- 11 quest, and

| 12 | if it is determined that the storage operating system is permitted to retrieve |
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| 13 | readahead data from the data container: |
| 14 | (i) selecting an amount of readahead data to retrieve from the |
| 15 | data container based on a plurality of factors stored within a readset data |
| 16 | structure associated with the read stream; and |
| 17 | (ii) retrieving the selected amount of readahead data from the |
| 18 | data container. |
| 1 | 23. (Original) The storage system of claim 22, wherein the data container is a file, direc- |
| 2 | tory, vdisk or lun. |
| 1 | 24. (Original) The storage system of claim 22, wherein the storage operating system is |
| 2 | determined to be permitted to retrieve readahead data from the data container when the |
| 3 | client-requested data extends the read stream past a predetermined next readahead value. |
| 1 | 25. (Original) The storage system of claim 24, wherein the predetermined next reada- |
| 2 | head value is updated based on a percentage of the selected amount of readahead data. |
| 1 | 26. (Previously Presented) The storage system of claim 22, wherein the plurality of fac- |
| 2 | tors used to select the amount of readahead data includes at least one of: |
| 3 | (i) the amount of client-requested data, |
| 4 | (ii) a number of client read requests processed in the read stream, and |
| 5 | (iii) a read-access style associated with the data container. |
| 1 | 27. (Original) The storage system of claim 22, wherein the selected amount of readahead |
| 2 | data is doubled if the number of client read requests processed in the read stream is |
| 3 | greater than a first threshold value. |

- 28. (Currently Amended) A computer-readable media comprising instructions for execu-
- tion in a processor for the practice of a method for a storage operating system imple-
- mented in a storage system to optimize the amount of readahead data retrieved for a read
- 4 stream established in a data container stored in the storage system, the method compris-
- 5 ing:

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- receiving a client read request at the storage system, the client read request indi-
- 7 cating client-requested data for the storage operating system to retrieve from the data
- 8 container containing the read stream;
 - determining whether the storage operating system is permitted to retrieve reada-
- head data from the data container in response to the received client read request;
 - if it is determined that the storage operating system is permitted to retrieve reada-
- head data from the data container, performing the steps of:
- (i) selecting an amount of readahead data to retrieve from the data container
- based on a plurality of factors stored within a readset data structure associ-
- ated with the read stream; and
 - (ii) retrieving the selected amount of readahead data from the data container.
- 29. (Original) The computer-readable media of claim 28, wherein the data container is a
- 2 file, directory, vdisk or lun.
- 30. (Previously Presented) The method of claim 1, wherein the retrieved readahead data
- 2 is stored in one or more buffers, the buffers containing a flush queue, the flush queue be-
- 3 ing configured to reuse buffers after a predetermined period of time.
- 1 31. (Previously Presented) The method of claim 30, wherein the read stream corresponds
- to a read-once data transfer and data retrieved from the data container is stored in the
- 3 flush queue.

- 32. (Previously Presented) The method of claim 30, wherein the retrieved readahead data
- 2 is stored in the flush queue.
- 33. (Previously Presented) The method of claim 30, wherein one or more buffers ac-
- 2 cessed from the flush queue are re-enqueued on a normal queue.
- 34. (Currently Amended) A method for optimizing readahead data retrieval for a read
- stream established in a data container stored in a storage system, the method comprising:
- receiving a client read request at the storage system, the client read request be-
- 4 longing to the read stream and indicating an amount of client-requested data;
- selecting an amount of readahead data based on the indicated amount of client-
- requested data stored within a readset data structure associated with the read stream; and
- retrieving the selected amount of readahead data from the data container.
- 35. (Previously Presented) The method of claim 34, wherein the selected amount of
- readahead data is set equal to a multiple of a predetermined amount, and wherein the
- multiple is associated with the amount of client-requested data.
- 1 36. (Previously Presented) The method of claim 34, wherein the selected amount of
- readahead data is set equal to a multiple of the amount of client-requested data.
- 37. (Previously Presented) The method of claim 36, further comprising the step of
- rounding the selected amount of readahead data to the size of a data block.
- 1 38. (Previously Presented) The method of claim 34, wherein the selected amount of
- readahead data is set equal to a predetermined upper limit.
- 39. (Currently Amended) A method for optimizing readahead data retrieval for a read
- stream established in a data container stored in a storage system, the method comprising:

| , | receiving a chefit read request at the storage system, the chefit read request be- |
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| 4 | longing to the read stream and indicating client-requested data; |
| 5 | selecting an amount of readahead data based on a read-access style associated |
| 6 | with the data container, wherein the read-access style is stored within a readset data struc |
| 7 | ture associated with the read stream; and |
| 8 | retrieving the selected amount of readahead data from the data container. |
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| 1 | 40. (Previously Presented) The method of claim 39, wherein the selected amount of |
| 2 | readahead data equals zero if the read-access style corresponds to a random read-access |
| 3 | style. |
| 1 | 41. (Previously Presented) A method for optimizing readahead data retrieval for a read |
| 2 | stream established in a data container stored in a storage system associated with a number |
| 3 | of storage devices, the method comprising: |
| 4 | receiving a client read request at the storage system, the client read request be- |
| 5 | longing to the read stream and indicating client-requested data; |
| 6 | selecting an amount of readahead data based on the number of storage devices; |
| 7 | and |
| 8 | retrieving the selected amount of readahead data from the data container. |
| 1 | 42. (Previously Presented) The method of claim 41, wherein the step of selecting an |
| 2 | amount of readahead data further comprises: |
| 3 | determining whether a flag is associated with the read stream, the flag indicating |
| 4 | that the storage system is associated with more than a predetermined number of storage |
| 5 | devices; and |
| 6 | in response to determining whether the flag is associated, selecting the amount of |
| 7 | readahead data. |

- 1 43. (Previously Presented) The method of claim 41, wherein the storage devices com-
- 2 prise one or more disks.
- 1 44. (Currently Amended) A method for optimizing readahead data retrieval for a read
- 2 stream established in a data container stored in a storage system, the method comprising:
- receiving a client read request at the storage system, the client read request be-
- 4 longing to the read stream and indicating client-requested data;
- selecting an amount of readahead data based on a plurality of factors stored within
- 6 a readset data structure associated with the read stream; and
- retrieving the selected amount of readahead data from the data container.
- 1 45. (Previously Presented) The method of claim 44, wherein the retrieved readahead
- data is stored in one or more buffers, the buffers containing a flush queue, the flush queue
- being configured to reuse buffers after a predetermined period of time.
- 46. (Previously Presented) The method of claim 45, wherein the read stream corre-
- sponds to a read-once data transfer and data retrieved from the data container is stored in
- 3 the flush queue.
- 47. (Previously Presented) The method of claim 45, wherein the retrieved readahead
- 2 data is stored in the flush queue.
- 48. (Previously Presented) The method of claim 45, wherein one or more buffers ac-
- 2 cessed from the flush queue are re-enqueued on a normal queue.
- 49. (Currently Amended) A system for optimizing readahead data retrieval for a read
- stream established in a data container stored in a storage system, the system comprising:
- means for receiving a client read request at the storage system, the client read re-
- 4 quest belonging to the read stream and indicating client-requested data;

- means for selecting an amount of readahead data based on a plurality of factors

 stored within a readset data structure associated with the read stream; and
- means for retrieving the selected amount of readahead data from the data con-
- 8 tainer.
- 50. (Previously Presented) The system of claim 49, wherein the retrieved readahead data
- 2 is stored in one or more buffers, the buffers containing a flush queue, the flush queue be-
- ing configured to reuse buffers after a predetermined period of time.
- 51. (Previously Presented) The system of claim 50, wherein the read stream corresponds
- to a read-once data transfer and data retrieved from the data container is stored in the
- 3 flush queue.
- 52. (Previously Presented) The system of claim 50, wherein the retrieved readahead data
- 2 is stored in the flush queue.
- 53. (Previously Presented) The system of claim 50, wherein one or more buffers ac-
- 2 cessed from the flush queue are re-enqueued on a normal queue.

- Please add new claims 54 et al.
- 1 54. (New) A method, comprising:
- receiving a plurality of client read requests at a storage system, the client read re-
- quests indicating client-requested data sets for a storage operating system to retrieve from
- 4 one or more data containers containing one or more read streams;
- selecting an amount of readahead data to retrieve from the one or more data con-
- tainers based on a plurality of factors stored within a readset data structure associated
- 7 with each read stream;
- retrieving the selected amount of readahead data from the data container;
- 9 processing one or more of the plurality of client read requests; and
- adjusting, as client requests are processed, the plurality of factors stored within
- the readset data structure associated with each read stream to optimize amount of reada-
- head data is cached for each read stream.
- 1 55. (New) The method of claim 54, further comprising:
- determining whether the storage operating system is permitted to retrieve reada-
- head data from the one or more data containers in response to each received client read
- 4 request.
- 56. (New) The method of claim 54, wherein the one or more data containers are at least
- one of a file, a directory, a vdisk or a lun.
- 1 57. (New) The method of claim 55, wherein the storage operating system is determined
- to be permitted to retrieve readahead data from the one or more data containers when the
- 3 client-requested data extends the read stream past a predetermined next readahead value.
- 1 58. (New) The method of claim 57, wherein the predetermined next readahead value is
- stored in a readset data structure associated with the read stream.

- 1 59. (New) The method of claim 57, wherein the predetermined next readahead value is
- 2 updated based on a percentage of the selected amount of readahead data.
- 1 60. (New) The method of claim 54, wherein a read-access style associated with the one
- or more data containers is one of the plurality of factors used to select the amount of
- 3 readahead data.
- 1 61. (New) The method of claim 60, wherein the selected amount of readahead data equals
- zero if the read-access style corresponds to a random read-access style.
- 62. (New) The method of claim 54, wherein a number of client read requests processed
- 2 in the read stream is one of the plurality of factors used to select the amount of readahead
- 3 data.
- 1 63. (New) The method of claim 62, wherein the number of client read requests processed
- 2 in the read stream is stored as a count value in a readset data structure associated with the
- 3 read stream.
- 1 64. (New) The method of claim 54, wherein the amount of client-requested data is one
- of the plurality of factors used to select the amount of readahead data.
- 1 65. (New) The method of claim 64, wherein the selected amount of readahead data is set
- equal to a predetermined upper limit for large amounts of client-requested data.
- 1 66. (New) The method of claim 54, wherein the selected amount of readahead data is
- doubled if the number of client read requests processed in the read stream is greater than
- 3 a first threshold value.

- 1 67. (New) The method of claim 55, wherein the client-requested data is identified as
- 2 read-once data when either (i) the number of client read requests processed in the read
- stream is greater than a second threshold value or (ii) a set of metadata associated with
- 4 the read stream indicates that the client-requested data is read-once data.
- 1 68. (New) The method of claim 54, wherein the selected amount of readahead data is
- stored in one or more buffers enqueued on a flush queue, the flush queue being config-
- 3 ured to reuse buffers after a predetermined period of time.